

### Overcoming Today's Limitations of Standard C++ with Kokkos

Damien Lebrun-Grandié

ORNL is managed by UT-Battelle LLC for the US Department of Energy



### Content

- Brief overview of Kokkos
- Memory and execution spaces
- Hierarchical parallelism







### The Kokkos ecosystem





### The Kokkos team







# Kokkos Core – Contributions and Usage



#### Usage

#### https://kokkosteam.slack.com >1000 Registered Users >130 Institutions





### What do we mean by "Standard C++"

- Code you can write today **without** using **language extensions** or **additional libraries**, which is portable to other compiler and systems and can be "automatically" accelerated with GPUs.
- Which excludes
  - CUDA/HIP require use of <u>host</u> and <u>device</u> attributes on functions and triple chevron syntax <<<...>>> for GPU kernel launches
  - OpenACC/OpenMP use #pragma directives to control GPU accelerations
  - Thrust or oneDPL let you express parallelism portably but only support a limited number of CPU and GPU backends





# Kokkos execution and memory spaces



# Standard C++ abstract machine

• Threads of execution evaluating functions that operate on objects that are in a flat storage space

• But...

- No notion of hierarchy (caches, etc.)
- No concept of host or device memory, nor accessibility



Adelstein Lelbach, CppCon 2018



### Revisions of the C++ standard





# GPU-enabled implementations today

#### NVIDIA HPC SDK

- Offloading of parallel algorithms to NVIDIA GPUs
- Enabled with the -stdpar[=gpu] option to NVC++
- Relies on CUDA Unified Memory for all data movement between CPU and GPU memory
- Automatically migrating data towards the processor using it

```
std::for_each(
    std::execution::par_unseq,
    v.begin(), v.end(),
    [](int& x) { x = x * x; });
```

#### Intel OneAPI DPC++

- Support for Intel, NVIDIA, and AMD GPU devices using oneDPL device execution policies (non-standard) std::par and std::par\_unseq run on the host via TBB or OpenMP backend
- Unified Shared Memory C-style memory management via sycl::malloc\_{device,host,shared}
- Or SYCL buffer objects to pass data to device (oneAPI runtime controls data movement)

#### std::fill(

oneapi::dpl::execution::make\_device\_policy(queue),
begin(v), end(v), 42);



### Target machine for Kokkos programming model





# Kokkos execution spaces

• Define where kernels get executed and what backend to use

(e.g. Serial, Threads, OpenMP, Cuda, HIP, ...)

• Execution space instances encapsulating CUDA/HIP stream or SYCL queue

Always available: DefaultExecutionSpace DefaultHostExecutionSpace





# Kokkos **memory** spaces

• Define "where" and "how" memory allocation and access take place

(e.g. HostSpace, SYCLDeviceUSMSpace, SYCLSharedUSMSpace, SYCLHostUSMSpace)

Always available: HostSpace SharedSpace SharedHostPinnedSpace ExecutionSpace::memory\_space ExecutionSpace::scratch\_memory\_space





# Accessibility of allocations from target devices

#### memory access violation?

#### Kokkos::parallel\_for(

"Fill", N, KOKKOS\_LAMBDA(int i) {v(i) = value; })

- Kokkos cannot introspect user-provided functors
- But it provides a facility to check accessibility and catch most bugs at compile-time rather than runtime

static\_assert(is\_accessible\_from<
 typename View::memory\_space,
 DefaultExecutionSpace>::value);







# Kokkos user remains in control of data placement

#### Explicit data movement

auto v\_host = create\_mirror\_view\_and\_copy(HostSpace(), v);

- Explicit data transfer between host and device
- no-op when view is already accessible from the CPU



#### Managed memory

• Managed memory accessible from all CPUs and GPUs in the system as a single, coherent memory image with a common address







# Hierarchical parallelism



# Example: inner product <y, A\*x>

```
double result = 0.;
for (int i = 0; i < N, ++i) {
    double Ax_i = 0.;
    for (int j = 0; j < M, ++j) {
        Ax_i += A(i, j) * x(j);
    }
    result += y(i) * Ax_i;
}
```



How to parallelize using C++17 parallel algorithms?



# Accelerate <y,A\*x> with standard parallelism

```
double result = std::reduce(
    std::execution::par_unseq,
    counting_iterator(0), counting_iterator(N),
    0.,
    [=](int i) {
        double Ax_i = 0.;
        for (int j = 0; j < M; ++j) {
            Ax_i += A(i, j) * x(j);
        }
        return y(i) * Ax_i;
    });</pre>
```



**Problem:** What if we don't have enough rows to saturate the GPU?



# Frontier compute node





[1x] 64-core AMD "Optimized 3rd Gen EPYC" CPU
[4x] MI250x each with 2 GCDs
Each GCD contains 110 CUs
64 GB of HBM accessible at 1.6 TB/s

CEs which dispatch wavefronts to CUs vavefronts from a single workgroup are gned to the same CU

Work items in a wavefront are scheduled in units of 64 called wavefronts Up to 64 KB of LDS can be allocated

Each CU has 4 MCUs and 4 16-wide SIMD units Each wavefront is assigned to a single 16-wide SIMD unit

Each CU maintains an instruction buffer for 10 wavefronts



# Hierarchical parallelism in Kokkos

- Exploit multiple level of shared-memory parallelism These levels include thread teams, threads within a team, and vector lanes.
- Able to nest these levels of parallelism, and execute parallel\_for(), parallel\_reduce(), or parallel\_scan() at each level
- Syntax differs only by the execution policy which is the 1<sup>st</sup> argument to the parallel\_\*
- Also exposing a "scratch pad" memory which provides thread private and team private allocations



# Accelerate <x,A\*y> with Kokkos

#### double result; parallel\_reduce( "yAx", TeamPolicy(N, AUTO), KOKKOS\_LAMBDA(auto const &team\_handle, double &partial\_result) { int const i = team\_handle.league\_rank(); double Ax i; parallel\_reduce( TeamThreadRange(team\_handle, M), [&](int const i, double & update) { update += A(i, j) \* x(j);}, Ax i); if (team\_handle.team\_rank() == 0) partial\_result += $y(i) * Ax_i$ ; result);





# **Sparse Matrix-Vector product (SpMV)**



The Kokkos EcoSystem: Comprehensive Performance Portability for High Performance Computing DOI: 10.1109/MCSE.2021.3098509 Kokkos Kernels native implementation with 3level hierarchical parallelism competes with vendor optimized libraries

rallel_for(	
"SpMV",	
TeamPolicy((nrows + rows_per_team - 1) / rows_per_team, team_size, 8), KOKKOS_LAMBDA(auto const& team) {	
<pre>int const first_row = team.league_rank() * rows_per_team; int const last_row = first_row + rows_per_team &lt; nrows</pre>	
? first_row + rows_per_team : nrows:	
parallel_for( 2nd [eve] TeamThreadRange(team, first_row, last_row), [&](int const row) { int const row_start = A.row_ptr(row); int const row_length = A.row_ptr(row + 1) - row_start; double y_row;	
<pre>parallel_reduce( 3rd level ThreadVectorRange(team, row_length), [=](int const i, double&amp; sum) {     sum +=         A.values(i + row_start) * x(A.col_idx(i + row_start));     },     y_row); </pre>	
y(row) = y_row; });	



### Wrap up

- We don't want memory management strategy to be dictated
- We want to be able to check accessibility of data
- We want nested algorithms



### Thank you!

# Damien L-G <lebrungrandt@ornl.gov>



Open slide master to edit